

Merge Sort $\rightarrow$ Divide & Conquer

$\rightarrow$ Good for LL
(Inplace)

Heap Sort

$\rightarrow$ clumsy for LL.
but good for arrays.

$\rightarrow$ Non-inplace for arrays.

Quick Sort $\rightarrow$ Best known comparison-based sorting

$\rightarrow$ D & C

$\text{SORT}(L, n)$

Base: $(n == 1) \Rightarrow L' = L$

Inductive:

Split

into two non-empty lists L_1 and L_2

Rec.

$L'_1 = \text{SORT}(L_1)$

$L'_2 = \text{SORT}(L_2)$

$L' = \text{MERGE}(L'_1, L'_2) \rightarrow \text{Rec.}$

Split in the Middle
 $\downarrow$
Merge Sort

Algorithm

$\text{mergeSort}(A, i, j)$

{

if $(i \geq j)$ return;

$\text{mid} = (i + j) / 2;$

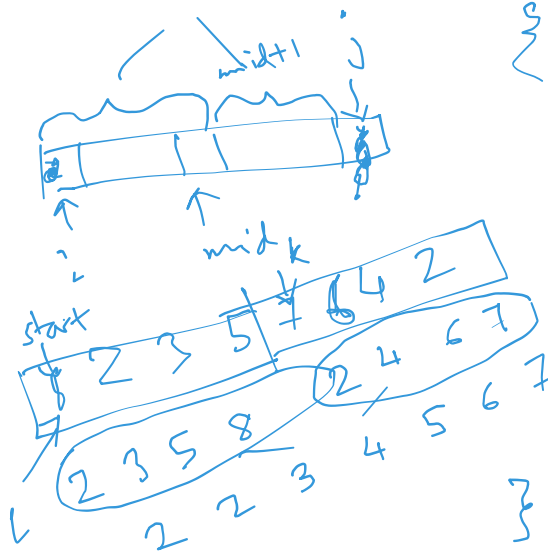
$\rightarrow \text{mergeSort}(A, i, \text{mid});$ -

$\rightarrow \text{mergeSort}(A, \text{mid} + 1, j);$ -

merge $(A, i, j);$ // Preserves the sorted order

}

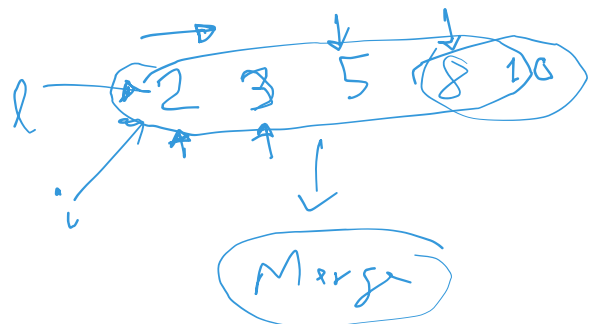
j



Merge:

$\text{merge}(A, i, j)$

{



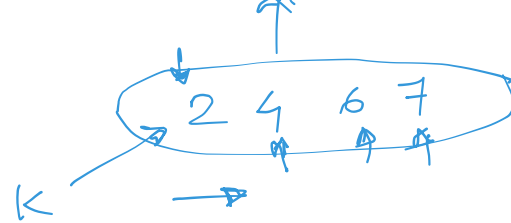
step

start = i;

mid = (i + j) / 2;

k = mid + 1;

l = i;



B → 2 2 3 4 5 6 7
8 10

merge

while (i ≤ mid & k ≤ j) {

if (A[i] ≤ A[k]) {

B[l++] = A[i++];

} else {

B[l++] = A[k++];

}

}

if (i > mid) {

for (; k ≤ j;) B[l++] = A[k++];

} else {

for (; i ≤ mid;) B[l++] = A[i++];

}

T(1) = 1

Exercise

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

0 1 3 4 5 7 8 9

output

3 5 7 9

0 1 4 8

3 7

5 9

4 8

0 1

7

3

9

5

8

4

1

0

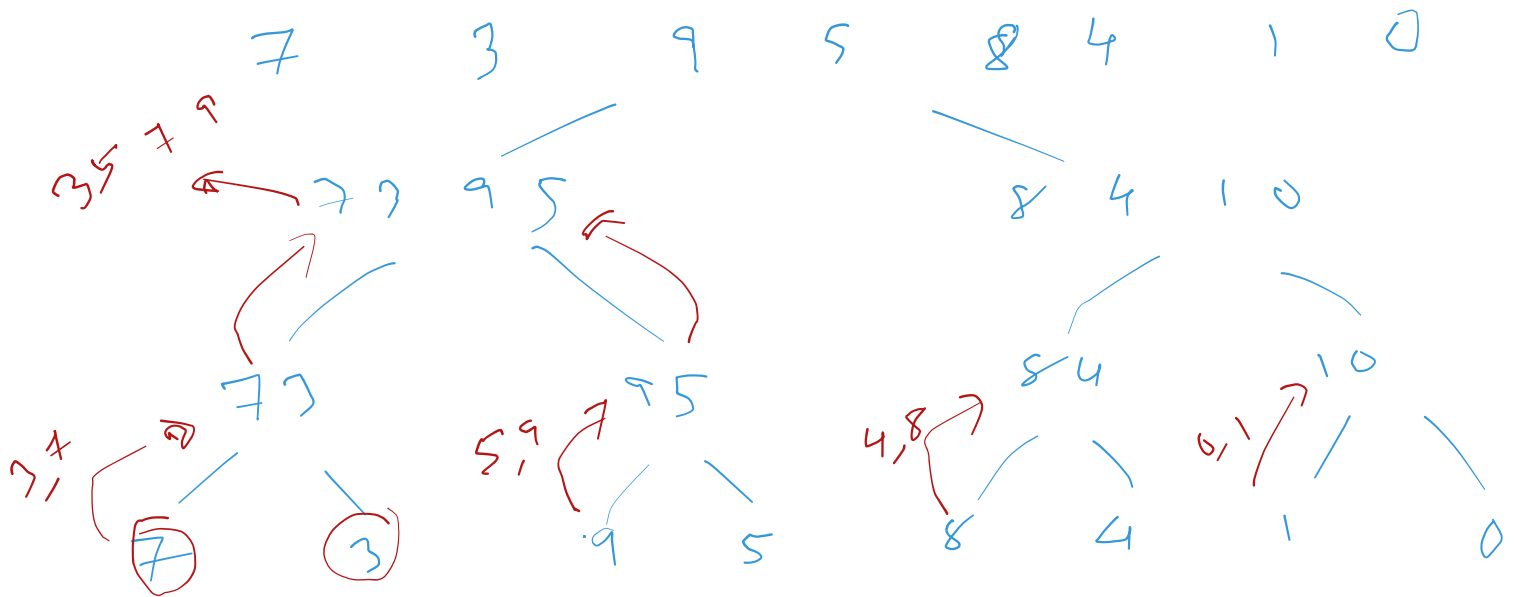
copy the remaining
entire

A[8]

Input

n

$\Theta(n \lg n)$



Natural sorting algo for Linked list
→ Merge sort

But Not an inplace sort for Arrays.