Logic

Readings: C&C, Ch. 4.1–4.3.1

1. Propositional logic

• *Propositional logic* is concerned with complex propositions built from simple propositions with the help of logical connectives.

1.1. Introducing logical connectives

- Assuming *p* and *q* are propositions:
	- conjunction: *p ∧ q*, equivalent to *p and q* in natural language;
	- disjunction: *p ∨ q*, equivalent to *p or q* in natural language;
	- material implication: $p \rightarrow q$, roughly equivalent to *if p, then q* in natural language;
	- negation: *¬p*, equivalent to *not p* in natural language.

1.2. Semantics of logical connectives via truth tables

1.2.1. Conjunction

- (1) All the figures have the same shape and all the figures have the same color.
- For [\(1\)](#page-0-0) to be true, it has to be the case that both conjuncts are true. That is the case for any sentence of the form *p ∧ q*. We can thus represent the semantics of conjunction using a *truth table*:

$$
\sf (2)
$$

1.2.2. Disjunction

- (3) All the figures have the same shape or all the figures have the same color.
- For [\(3\)](#page-0-1) to be true, one or both disjuncts have to be true. That is the case for any sentence of the form $p \vee q$.
- You might have an intuition that a sentence of the form *p or q* is true only if one of the disjuncts is true, but not both. This reading corresponds to *exclusive* (as opposed to *inclusive*) *disjunction*. In propositional logic we sometimes write *p xor q* to represent exclusive disjunction.

• Draw the truth tables for *p ∨ q* and *p xor q*.

• Exclusive disjunction seems to be the default reading of English sentences involving *or*. It is commonly assumed, however, that this reading arises as a conversational implicature. That is, *or* is taken to denote inclusive disjunction and this meaning is strengthened via Gricean reasoning. Go through the steps of the Gricean reasoning process (i.e., which maxims are involved etc.).

• Show that the exclusive reading is indeed a conversational implicature.

1.2.3. Material implication

(4) *Context: You are playing a slot machine that displays three figures when you pull the lever. You win if and only if all the figures have the same shape.*

If all the figures have the same shape, you win.

- Given the context, the sentence in [\(4\)](#page-1-0) should be true if the slot machine is working well. In order to check whether it really does work well, one could go through the following scenarios:
	- All the figures have the same shape and you win the machine is working well.
	- The figures don't all have the same shape and you win the machine is working well.
	- All the figures have the same shape and you don't win problem!
	- The figures don't all have the same shape and you don't win the machine is working well.
- In other words, [\(4\)](#page-1-0) is only false when its *antecedent* is true and its *consequent* is false. We will assume that the same holds for any sentence of the form $p \rightarrow q$.

• Draw the truth table for $p \rightarrow q$.

- To capture the meaning of natural language conditionals (*if p, (then) q*), we often need something more sophisticated then material implication. E.g., treating [\(5\)](#page-2-0) via material implication would predict it to be true, but is [\(5\)](#page-2-0) intuitively true? We will not talk about natural language conditionals in this course, however.
- (5) If the moon is made of green cheese, then I had yogurt for breakfast this morning.

1.2.4. Negation

- The connectives above are all *binary connectives*. Negation is a *unary connective*.
- (6) It's not the case that all the figures have the same shape.
- [\(6\)](#page-2-1) is true if and only if the sentence *All the figures have the same shape* is false. More generally: $\neg p$ is true iff *p* is false.

In-class Exercise 3

• Draw the truth table for $\neg p$.

1.3. Reasoning with truth tables

- Reminder: φ ([fai]) entails ψ ([sai]) iff whenever φ is true, ψ is also true.
- We can thus see if *φ* entails *ψ* by looking at all the rows of the truth table where *φ* is true and checking if ψ is also true in those rows.

In-class Exercise 4

- Does *p* entail *p ∨ q*? Explain using the relevant truth table from In-class Exercise 1.
- Does $[p \rightarrow q] \land \neg p$ entail $\neg q$? Explain using a truth table.

- φ contradicts ψ iff whenever φ is true, ψ is false (and vice versa).
- We can thus see if *φ* contradicts *ψ* by looking at the truth table and checking if in all rows the values of φ and ψ are different.

In-class Exercise 5

- Show that *p* contradicts $\neg p$ using the truth table from In-class Exercise 3.
- Does $p \rightarrow q$ contradict $p \land \neg q$? Explain using a truth table.

1.4. Grammar of propositional logic

• Propositional logic is a language, so it has a syntax (specifying which formulas are well-formed) and a semantics (assigning denotations to all well-formed formulas).

1.4.1. Syntax of propositional logic

- Syntactic rules of propositional logic:
	- 1. Atomic formulas: *p*, *q*, *r*, ...
	- 2. Negation: if φ is a formula, then $\neg \varphi$ is a formula.
	- 3. Binary connectives: if φ and ψ are both formulas, then so are:
		- (a) [*φ ∧ ψ*]
		- (b) [*φ ∨ ψ*]
		- (c) [*φ → ψ*]
		- (d) $[\varphi \leftrightarrow \psi]$
	- 4. Nothing else is a formula.
- Brackets are important. E.g., [*¬p ∧ q*] and *¬*[*p ∧ q*] are two different formulas. Outermost brackets are usually omitted.

In-class Exercise 6

- For each of the strings below say if it is a formula of propositional logic.
- (7) a. *¬¬p* b. $\neg \wedge p$ c. $p, q \vee r$

1.4.2. Semantics of propositional logic

- The truth of a formula depends on what the world is like, i.e., on the *model M*.
- In propositional logic a model will specify for each simple proposition whether it's true or false. For example, if our model M_1 specifies that p is true and q is false, we can write:
- $M_1(p) = 1$ $M_1(q) = 0$
- *Semantic values* are assigned to linguistic expressions by the *valuation function*, written as $\lceil \cdot \rceil$, which will be relativized to a model: $[\![\cdot]\!]^M$.
- Semantics of propositional logic:
	- 1. Atomic formulas: if φ is an atomic formula, then $[\![\varphi]\!]^M = M(\varphi)$ (i.e., whichever truth value the model *M* specifies for φ).
	- 2. Negation: $[\![\neg \varphi]\!]^M = \mathrm{I} \text{ iff } [\![\varphi]\!]^M = \circ.$
	- 3. Binary connectives:
		- $\left[\varphi \wedge \psi \right]^M = 1$ iff $\left[\varphi \right]^M = 1$ and $\left[\psi \right]^M = 1$
		- (b) $[\varphi \vee \psi]^{M} = 1$ iff $[\varphi]^{M} = 1$ or $[\psi]^{M} = 1$
		- ϕ (\mathbf{c}) $[\varphi \to \psi]^{M} = 1$ iff $[\![\varphi]\!]^{M} = 0$ or $[\![\psi]\!]^{M} = 1$
		- (d) $[\![\varphi \leftrightarrow \psi]\!]^M = 1$ iff $[\![\varphi]\!]^M = [\![\psi]\!]^M$

• Assuming a model M_1 such that $[\![p]\!]^{M_1} = 1$ and $[\![q]\!]^{M_1} = 0$, compute the semantic value of the formulas below step by step. For each step say which semantic rule of propositional logic you used.

$$
(9) \quad \text{a.} \quad \llbracket \neg [p \wedge q] \rrbracket^{M_1} =
$$

$$
\mathbf{b.} \quad [\![\neg p \vee q]\!]^{M_1} =
$$

$$
\mathbf{c.} \quad [p \to \neg q]^{M_1} =
$$

$$
d. \quad [\neg p \leftrightarrow q]^{M_1} =
$$

2. Predicate logic

• Predicate logic (a.k.a *first-order logic*, *first-order predicate calculus*) adds predication and quantification to propositional logic.

2.1. Predication

- In propositional logic, we'd express the natural language sentences in (10) as atomic sentences *p*, *q* and *r*. But this can't capture what these sentences have in common (namely, (10a) and (10b) have the same subject, and (10b) and (10c) the same predicate).
	- (10) a. Neil giggles.
		- b. Neil sings.
		- c. Marilyn sings.
- In predicate logic, we have *individual constants*, e.g., n (Neil) and m (Marilyn), and we have *predicates*, e.g., Giggles and Sings. This allows us to express the sentences in (10) as in (11):
	- (n) a. Giggles (n)
		- b. Sings(n)
		- c. Sings(m)

• In addition to one-place predicates like Sings, there are also two-place predicates like Sees, and three-place predicates like Gives.

 (I_2) a. Sees (m,n) b. Gives(n,m,e)

- Just as in propositional logic, atomic sentences can be connected with *∧*, *∨*, *→* and *↔* to form complex sentences.
	- $(i3)$ Sings(n) \rightarrow Giggles(m)

In-class Exercise 8

- For each of the strings below say if it is a formula of predicate logic.
	- (14) a. *¬¬*Purrs(c)
		- b. Smart(Student)
		- c. Student(h,r)
		- d. Kissed(r)
		- e. Petted(h,c) *∨* Petted(c,h)
		- f. Likes $(h,r) \leftrightarrow$ Likes (r,h)

2.2. Quantification

- • Predicate logic also has *variables* (as opposed to *constants*) and quantifiers. With these tools, we can capture meanings of sentences such as:
	- (15) a. Some cat purrs.
		- b. Every cat purrs.
- • [\(15a\)](#page-6-0) is true iff there is an individual *x* such that *x* is a cat and *x* purrs. To represent this meaning, we will use the symbol *∃*, which stands for the *existential quantifier* (meaning 'there exists' or 'there is'), and variables:

(16) *∃x.*[Cat(*x*) *∧* Purrs(*x*)]

• [\(15b\)](#page-6-1) is true iff for all *x*, if *x* is a cat, it purrs. To represent this meaning, we will use the symbol *∀*, which stands for the *universal quantifier* (meaning 'for all'), and variables:

 (I7) $\forall x.$ [Cat $(x) \rightarrow$ Purrs (x)]

- • Brackets are used to indicate the *scope* of the quantifier (but are often omitted when the scope is unambiguous). A quantifier only *binds* the variables that are in its scope. E.g., in [\(18a\)](#page-6-2) *∃* binds both occurrences of *x*, but in [\(18b\)](#page-6-3) the second occurrence of *x* is *unbound*.
	- (18) a. $\exists x.$ [Cat $(x) \wedge$ Purrs (x)] b. *∃x.*[Cat(*x*)] *∧* Purrs(*x*)

- Read the following formulas aloud using such words as *there exists*, *for all*, *it's not the case that*, etc.:
	- (19) a. *∀x.*[*∃y.*[Likes(*x, y*)]] b. *∀x.*[Student(*x*) *→ ∃y.*[Cat(*y*) *∧* Petted(*x, y*)]] c. *¬∃x.*[Cat(*x*) *∧* Purrs(*x*)] d. $\forall x. \neg [Cat(x) \rightarrow Purr(x)]$
- Express the following sentences in predicate logic (you can use as many constants as you need):
	- (20) a. Some student is tall.
		- b. Some man petted some cat.
		- c. It's not the case that some woman petted Fido.
		- d. All cats are white or gray.

2.3. Semantics of predicate logic

- The semantics of predicate logic is a bit more involved than that of propositional logic and we will skip the details here (if you are interested, you can have a look at Coppock and Champollion textbook draft, Ch. 3.2–3).
- But on an intuitive level, if we understand what a sentence in predicate logic expresses, we can also determine which semantic value this sentence receives in a given model.

In-class Exercise 10

• Assume a model M_2 , where the set of individuals consists of Chloe and Rocky ($\{c, r\}$), who are both cats. Only Chloe purrs. Chloe scratches Rocky but Rocky doesn't scratch Chloe. Determine which semantic value the sentences below receive in *M*2.

$$
(2I) \qquad a. \qquad Purrs(r)
$$

b.
$$
\forall x. [Cat(x)]
$$

- c. *∃x.*[Cat(*x*) *∧* Purrs(*x*)]
- d. $∀x.∀y.$ [Scratches(*x*, *y*) → Purrs(*x*)]
- e. *∀x.∀y.*[Scratches(*x, y*) *↔* Scratches(*y, x*)]

What you need to know

Key notions: exclusive vs. inclusive disjunction, antecedent, consequent, binary vs. unary connectives, model, valuation function, individual constant, variable, existential quantifier, universal quantifier, quantifier scope, bound variable, unbound variable

Skills:

- Draw truth tables for complex sentences in propositional logic.
- Determine whether two complex sentences in propositional logic entail or contradict one another using truth tables.
- Say if a given string is a well-formed formula in propositional logic.
- Read formulas written in propositional and/or predicate logic using such words as *there exists*, *for all*, *it's not the case that*, etc.
- Represent meanings of natural language sentences consisting of names, predicates, quantifiers such as *some cat* and *every cat*, and counterparts of logical connectives (*and*, *or*, *if... then*, and negation) using predicate logic.
- Compute semantic values of formulas in propositional logic given a model.